Fortran 90 Tutorial

Michael Metcalf

CN Division, CERN, CH 1211, Geneva 23, Switzerland

Full details of all the items in this tutorial can be found in *Fortran 90 Explained,* by M. Metcalf and J. Reid, (Oxford, 1990), the book upon which it has been based.

Fortran 90 contains the whole of FORTRAN 77—only the new features are described in this tutorial.

The tutorial is also available on WWW using the URL
`http://asis01.cern.ch/CN/CNTUT/f90/Overview.html`.

The author wishes to thank Michel Goossens (CERN/CN) for his helpful and skilful assistance in preparing this tutorial.

1. Language Elements

The basic components of the Fortran language are its *character set*. The members are:

- the letters `A ... Z` and `a ... z` (which are equivalent outside a character context);
- the numerals `0 ... 9`;
- the underscore `_` and
- the special characters

```
=  :  +  blank  -  *  /  (  )  ,  .  $  ' (old)
!  "  %  &  ;   <  >  ?                  (new)
```

From these components, we build the *tokens* that have a syntactic meaning to the compiler. There are six classes of token:

Label: `123`
Constant: `123.456789_long`
Keyword: `ALLOCATABLE`
Operator: `.add.`
Name: `solve_equation` (can have up to 31 characters, including a `_`).
Separator: `/ ( ) (/ /) , = => :  ::  ; %`

From the tokens, we can build statements. These can be coded using the new free *source form* which does not require positioning in a rigid column structure, as follows:

```
FUNCTION string_concat(s1, s2)     ! This is a comment
   TYPE (string), INTENT(IN) :: s1, s2
   TYPE (string) string_concat
   string_concat%string_data = s1%string_data(1:s1%length) // &
      s2%string_data(1:s2%length)  ! This is a continuation
   string_concat%length = s1%length + s2%length
END FUNCTION string_concat
```

Note the trailing comments and the trailing continuation mark. There may be 39 continuation lines, and 132 characters per line. Blanks are significant. Where a token or character constant is split across two lines:

```
        ...        start_of&
 &_name
        ...   'a very long &
 &string'
```

a leading `&` on the continued line is also required.

Automatic conversion of source form for existing programs can be carried out by `CONVERT` (CERN Program Library Q904). Its options are:

- significant blank handling;
- indentation;
- `CONTINUE` replaced by `END DO`;
- name added to subprogram `END` statement; and
- `INTEGER*2` etc. syntax converted.

The source code of the `CONVERT` program can be obtained by anonymous ftp to `jkr.cc.rl.ac.uk` (130.246.8.23). The directory is `/pub/MandR` and the file name is `convert.f90`.

Fortran has five *intrinsic data types*. For each there is a corresponding form of *literal constant*. For the three numeric intrinsic types they are:

`INTEGER   1   0   -999   32767   +10`

for the default *kind*; but we may also define, for instance for a desired range of -10^4 to 10^4, a named constant, say `two_bytes`:

```
INTEGER, PARAMETER :: two_bytes &
                    = SELECTED_INT_KIND(4)
```

that allows us to define constants of the form

```
-1234_two_bytes
+1_two_bytes
```

Here, `two_bytes` is the kind type parameter; it can also be a default integer literal constant, like

```
-1234_2
```

but use of an explicit literal constant would be non-portable.

The `KIND` function supplies the value of a kind type parameter:

```
KIND(1)
KIND(1_two_bytes)
```

and the `RANGE` function supplies the actual decimal range (so the user must make the actual mapping to bytes):

```
RANGE(1_two_bytes)
```

Also, in `DATA` statements, binary, octal and hexadecimal constants may be used:

```
B'01010101'
O'01234567'
Z'10fa'
```

`REAL`

There are at least two real kinds – the default, and one with greater precision (this replaces `DOUBLE PRECISION`). We might specify

```
INTEGER, PARAMETER :: long &
                = SELECTED_REAL_KIND(9, 99)
```

for at least 9 decimal digits of precision and a range of 10^{-99} to 10^{99}, allowing

```
1.7_long
```

Also, we have the intrinsic functions

```
KIND(1.7_long)
PRECISION(1.7_long)
RANGE(1.7_long)
```

that give in turn the kind type value, the actual precision (here at least 9), and the actual range (here at least 99).

`COMPLEX`

This data type is built of two integer or real components:

```
(1, 3.7_long)
```

The numeric types are based on model numbers with associated inquiry functions (whose values are independent of the values of their arguments):

`DIGITS(X)`	Number of significant digits
`EPSILON(X)`	Almost negligible compared to one (real)
`HUGE(X)`	Largest number
`MAXEXPONENT(X)`	Maximum model exponent (real)
`MINEXPONENT(X)`	Minimum model exponent (real)
`PRECISION(X)`	Decimal precision (real, complex)
`RADIX(X)`	Base of the model
`RANGE(X)`	Decimal exponent range
`TINY(X)`	Smallest postive number (real)

and these are important for portable numerical software.

The forms of literal constants for the two non-numeric data types are:

```
CHARACTER  'A string'
       "Another"
       'A "quote"'   ''
```

(the last being a null string). Other kinds are allowed, especially for support of non-European languages:

```
2_'   '
```

and again the kind value is given by the `KIND` function:

```
KIND('ASCII')
```

`LOGICAL`

Here, there may also be different kinds (to allow for packing into bits):

```
.FALSE.
.true._one_bit
```

and the `KIND` function operates as expected:

```
KIND(.TRUE.)
```

We can specify scalar *variables* corresponding to the five intrinsic types:

```
INTEGER(KIND=2) i
REAL(KIND=long) a
COMPLEX         current
LOGICAL         Pravda
CHARACTER(LEN=20) word
CHARACTER(LEN=2, KIND=Kanji) kanji_word
```

where the optional `KIND` parameter specifies a non-default kind, and the `LEN=` specifier replaces the `*len` form. The explicit `KIND` and `LEN` specifiers are optional and the following works just as well:

```
CHARACTER(2, Kanji) kanji_word
```

For *derived-data* types we must first define the form of the type:

```
TYPE person
   CHARACTER(10) name
   REAL          age
END TYPE person
```

and then create structures of that type:

```
TYPE(person) you, me
```

To select components of a derived type, we use the `%` qualifier:

```
you%age
```

and the form of a literal constant of a derived type is shown by:

```
you = person('Smith', 23.5)
```

which is known as a structure constructor.

Definitions may refer to a previously defined type:

```
TYPE point
   REAL x, y
END TYPE point
TYPE triangle
   TYPE(point) a, b, c
END TYPE triangle
```

and for a variable of type `triangle`, as in

```
TYPE(triangle) t
```

we then have components of type `point`:

```
t%a   t%b   t%c
```

which, in turn, have ultimate components of type `real`:

```
t%a%x   t%a%y   t%b%x   etc.
```

We note that the `%` qualifier was chosen rather than `.` because of ambiguity difficulties.

Arrays are considered to be variables in their own right. Given

```
REAL A(10)
INTEGER, DIMENSION(0:100, -50:50) :: map
```

(the latter an example of the syntax that allows grouping of attributes to the left of `::` and of variables sharing those attributes to the right), we have two arrays whose elements are in array element order (column major), but not necessarily in contiguous storage. Elements are, for example,

```
a(1)               a(i*j)
```

and are scalars. The subscripts may be any scalar integer expression. Sections are

```
a(i:j)               ! rank one
map(i:j, k:l:m)      ! rank two
a(map(i, k:l))       ! vector subscript
a(3:2)               ! zero length
```

Whole arrays and array sections are array-valued objects. Array-valued constants (constructors) are available:

```
(/ 1, 2, 3, 4, 5 /)
(/ (i, i = 1, 9, 2) /)
(/ ( (/ 1, 2, 3 /), i = 1, 10) /)
(/ (0, i = 1, 100) /)
(/ (0.1*i, i = 1, 10) /)
```

making use of the implied-`DO` loop notation familiar from I/O lists. A derived data type may, of course, contain array components:

```
TYPE triplet
   REAL, DIMENSION(3) :: vertex
END TYPE triplet
TYPE(triplet), DIMENSION(1) :: t
```

so that

```
t(2)                 ! a scalar (a structure)
t(2)%vertex          ! an array component of a scalar
```

There are some other interesting character extensions. Just as a substring as in

```
CHARACTER(80), DIMENSION(60) :: page
... = page(j)(i:i)          ! substring
```

was already possible, so new are the substrings

```
'0123456789'(i:i)
you%name(1:2)
```

Also, zero-length strings are allowed:

```
page(j)(i:i-1)       ! zero-length string
```

Finally, there are some new intrinsic character functions:

```
ACHAR                 IACHAR  (for ASCII set)
ADJUSTL               ADJUSTR
LEN_TRIM              INDEX(s1, s2, BACK=.TRUE.)
REPEAT                SCAN  (for one of a set)
TRIM                  VERIFY(for all of a set)
```

2. Expressions and Assignments

The rules for *scalar numeric* expresions and assignments, as known from FORTRAN 77, are extended to accommodate the non-default kinds we encountered in chapter 1. Thus, the mixed-mode numeric expression and assignment rules incorporate different kind type parameters in an expected way:

```
real2 = integer + real1
```

converts integer to a real value of the same kind as real1; the result is of same kind, and is converted to the kind of real2 for assignment.

For *scalar relational* operations, there is a set of new, alternative operators:

```
<    <=    ==    /=    >    >=
```

so we can write expressions such as

```
IF (a < b .AND. i /= j) THEN ! for numeric variables
flag = flag == semaphore     ! for logical variables
```

In the case of *scalar characters*, two old restrictions are lifted. Given

```
CHARACTER(8) result
```

it is now legal to write

```
result(3:5) = result(1:3)    ! overlap allowed
result(3:3) = result(3:2)    ! no assignment of null string
```

For an operation between derived-data types, or between a derived type and an intrinsic type, we must define the meaning of the operator. (Between intrinsic types, there are intrinsic operations only.) Given

```
TYPE string
   INTEGER       length
   CHARACTER(80) value
END TYPE string
CHARACTER    char1, char2, char3
TYPE(string) str1,  str2,  str3
```

we can write

```
str3  = str1//str2       ! must define operation
str3  = str1.concat.str2 ! must dedine operation
char3 = char2//char3     ! intrinsic operator only
str3  = char1            ! must define assignment
```

For the first three cases, assignment applies on a component-by-component basis (but can be overridden), and the first two cases require us to define the exact meaning of the // symbol. We see here the use both of an intrinsic symbol and of a named operator, `.concat.` . A difference is that, for an intrinsic operator token, the usual precedence rules apply, whereas for named operators their precedence is the highest as a unary operator or the lowest as a binary one. In

```
vector3 = matrix    *    vector1  + vector2
vector3 =(matrix .times. vector1) + vector2
```

the two expresions are equivalent only if appropriate parentheses are added as shown. In each case, we have to provide, in a module, procedures defining the operator and assignment, and make the association by an interface block, also in the module (we shall return to this later).

For the moment, here is an example of an interface for string concatenation

```
INTERFACE OPERATOR(//)
   MODULE PROCEDURE string_concat
END INTERFACE
```

and an example of part of a module containing the definitions of character-to-string and string to character assignment. The string concatenation function was shown already in chapter 1.

```
MODULE string_type
   TYPE string
      INTEGER length
      CHARACTER(LEN=80)   :: string_data
   END TYPE string
   INTERFACE ASSIGNMENT(=)
      MODULE PROCEDURE c_to_s_assign, s_to_c_assign
   END INTERFACE
   INTERFACE OPERATOR(//)
      MODULE PROCEDURE string_concat
   END INTERFACE
CONTAINS
```

```
  SUBROUTINE c_to_s_assign(s, c)
    TYPE (string), INTENT(OUT)    :: s
    CHARACTER(LEN=*), INTENT(IN)  :: c
    s%string_data = c
    s%length = LEN(c)
  END SUBROUTINE c_to_s_assign
  SUBROUTINE s_to_c_assign(c, s)
    TYPE (string), INTENT(IN)     :: s
    CHARACTER(LEN=*), INTENT(OUT) :: c
    c = s%string_data(1:s%length)
    END SUBROUTINE s_to_c_assign
    FUNCTION string_concat(s1, s2)
        :
  END FUNCTION string_concat
END MODULE string_type
```

Defined operators such as these are required for the expressions that are allowed too in structure constructors (see chapter 1):

```
str1 = string(2, char1//char2)  ! structure constructor
```

So far we have discussed scalar variables. In the case of *arrays*, as long as they are of the same shape (conformable), operations and assignments are extended in an obvious way, on an element-by-element basis. For

```
REAL, DIMENSION(10, 20) :: a, b, c
REAL, DIMENSION(5)      :: v, w
LOGICAL                    flag(10, 20)
```

can write

```
a = b                    ! whole array assignment
c = a/b                  ! whole array division and assignment
c = 0.                   ! whole array assignment of scalar value
w = v + 1.               ! whole array addition to scalar value
w = 5/v + a(1:5, 5)      ! array division, and addition to section
flag = a==b              ! whole array relational test and assignment
c(1:8, 5:10) = a(2:9, 5:10) + b(1:8, 15:20)
                         ! array section addition and assignment
v(2:5) = v(1:4)          ! overlapping section assignment
```

The order of expression evaluation is not specified in order to allow for optimization on parallel and vector machines. Of course, any operators for arrays of derived type must be defined.

There are some new real intrinsic functions that are useful for numeric computations:

```
CEILING        FLOOR          MODULO (also integer)
EXPONENT       FRACTION
NEAREST        RRSPACING      SPACING
SCALE          SET_EXPONENT
```

Like all FORTRAN 77 functions (`SIN`, `ABS`, etc., but not `LEN`), these are array valued for array arguments (i.e. are elemental).

3. Control Statements

The CASE construct is a replcement for the computed GOTO, but is better structured and does not require the use of statement labels:

```
SELECT CASE (number)       ! NUMBER of type integer
CASE (:-1)                 ! all values below 0
   n_sign = -1
CASE (0)                   ! only 0
   n_sign = 0
CASE (1:)                  ! all values above 0
   n_sign = 1
END SELECT
```

Each CASE selector list may contain a list and/or range of integers, character or logical constants, whose values may not overlap within or between selectors:

```
CASE (1, 2, 7, 10:17, 23)
```

A default is available:

```
CASE DEFAULT
```

There is only one evaluation, and only one match.

A simplified but sufficient form of the DO construct is illustrated by

```
outer: DO
inner:    DO i = j, k, l      ! only integers
             :
             IF (...) CYCLE
             :
             IF (...) EXIT outer
          END DO inner
       END DO outer
```

where we note that loops may be named so that the EXIT and CYCLE statements may specify which loop is meant.

Many, but not all, simple loops can be replaced by array expressions and assignments, or by new intrinsic functions. For instance

```
tot = 0.
DO i = m, n
   tot = tot + a(i)
END DO
```

becomes simply

```
tot = SUM( a(m:n) )
```

4. Program Units and Procedures

In order to discuss this topic we need some definitions. In logical terms, an executable program consists of one *main program* and zero or more *subprograms* (or *procedures*) - these do something. Subprograms are either *functions* or *subroutines*, which are either *external, internal* or *module* subroutines. (External subroutines are what we know from FORTRAN 77.)

From an organizational point of view, however, a complete program consists of *program units*. These are either *main programs, external subprograms* or *modules* and can be separately compiled.

An internal subprogram is one *contained* in another (at a maximum of one level of nesting) and provides a replacement for the statement function:

```
SUBROUTINE outer
   REAL x, y
   :
CONTAINS
   SUBROUTINE inner
      REAL y
      y = x + 1.
      :
   END SUBROUTINE inner     ! SUBROUTINE mandatory
END SUBROUTINE outer
```

We say that outer is the *host* of inner, and that inner obtains access to entities in outer by *host association* (e.g. to x), whereas y is a *local* variable to inner. The *scope* of a named entity is a *scoping unit*, here outer less inner, and inner.

The names of program units and external procedures are *global*, and the names of implied-DO variables have a scope of the statement that contains them.

Modules are used to package

- global data (replaces COMMON and BLOCK DATA);
- type definitions (themselves a scoping unit);
- subprograms (which among other things replaces the use of ENTRY);
- interface blocks (another scoping unit, see next article);
- namelist groups.

An example of a module containing a type defition, interface block and function subprogram is:

```
MODULE interval_arithmetic
   TYPE interval
      REAL lower, upper
   END TYPE interval
   INTERFACE OPERATOR(+)
      MODULE PROCEDURE add_intervals
   END INTERFACE
   :
CONTAINS
   FUNCTION add_intervals(a,b)
      TYPE(interval), INTENT(IN) :: a, b
      TYPE(interval) add_intervals
      add_intervals%lower = a%lower + b%lower
      add_intervals%upper = a%upper + b%upper
   END FUNCTION add_intervals             ! FUNCTION mandatory
   :
END MODULE interval_arithmetic
```

and the simple statement

```
USE interval_arithmetic
```

provides *use association* to all the module's entities. Module subprograms may, in turn, contain internal subprograms.

Arguments

We may specify the *intent* of dummy arguments:

```
     SUBROUTINE shuffle (ncards, cards)
        INTEGER, INTENT(IN)  :: ncards                    ! input values
        INTEGER, INTENT(OUT), DIMENSION(ncards) :: cards  ! output values
```

Also, INOUT is possible: here the actual argument must be a variable (unlike the default case where it may be a constant). Arguments may be *optional*:

```
     SUBROUTINE mincon(n, f, x, upper, lower, equalities, inequalities, convex, xstart)
        REAL, OPTIONAL, DIMENSION :: upper, lower
        .
        .
```

allows us to call `mincon` by

```
        CALL mincon (n, f, x, upper)
```

and in `mincon` we have someting like:

```
        IF (PRESENT(lower)) THEN   ! test for presence of actual argument
```

Arguments may be *keyword* rather than positional (which come first):

```
        CALL mincon(n, f, x, equalities=0, xstart=x0)
```

Optional and keyword arguments are handled by interface blocks.

Interface blocks

Any reference to an internal or module subprogram is through an interface that is "explicit" (that is, the compiler can see all the details). A reference to an external (or dummy) procedure is usually "implicit" (the compiler assumes the details). However, we can provide an explicit interface in this case too. It is a copy of the header, specifications and `END` statement of the procedure concerned, either placed in a module or inserted directly:

```
     REAL FUNCTION minimum(a, b, func)
! returns the minimum value of the function func(x) in the interval (a,b)
        REAL, INTENT(in) :: a, b
        INTERFACE
           REAL FUNCTION func(x)
              REAL, INTENT(IN) :: x
           END FUNCTION func
        END INTERFACE
        REAL f,x
        :
        f = func(x)   ! invocation of the user function.
        :
     END FUNCTION minimum
```

An interface block is obligatory for: optional and keyword arguments, `POINTER` and `TARGET` arguments (see later article), a `POINTER` function result (later) and new-style array arguments and array functions (later). Placed in a module, it allows full checks at compile time between actual and dummy arguments.

Overloading and generic interfaces

Interface blocks provide the mechanism by which we are able to define generic names for specific procedures:

```
INTERFACE gamma                    ! generic name
   FUNCTION sgamma(X)              ! specific name for low precision
      REAL (SELECTED_REAL_KIND( 6)) sgamma, x
   END
   FUNCTION DGAMMA(X)              ! specific name for high precision
      REAL (SELECTED_REAL_KIND(12)) dgamma, x
   END
END INTERFACE
```

where a given set of specific names corresponding to a generic name must all be of functions or all of subroutines.
We can use existing names, e.g. SIN, and the compiler sorts out the correct association.
We have already seen the use of interface blocks for defined operators and assignment (see Part 2).

Recursion

Indirect recursion is useful for multi-dimensional integration. To calculate

```
volume = integrate(fy, ybounds)
```

we might have

```
RECURSIVE FUNCTION integrate(f, bounds)
   ! Integrate f(x) from bounds(1) to bounds(2)
   REAL integrate
   INTERFACE
      FUNCTION f(x)
         REAL f, x
      END FUNCTION f
   END INTERFACE
   REAL, DIMENSION(2), INTENT(IN) :: bounds
   :
END FUNCTION integrate
```

and to integrate f(x, y) over a rectangle

```
FUNCTION fy(y)
   USE func            ! module func contains function f
   REAL fy, y
   yval = y
   fy = integrate(f, xbounds)
END
```

Direct recursion is when a procedure calls itself, as in

```
RECURSIVE FUNCTION factorial(n) RESULT(res)
   INTEGER res, n
   IF(n.EQ.1) THEN
      res = 1
   ELSE
      res = n*factorial(n-1)
   END IF
END
```

Here, we note the `RESULT` clause and termination test.

5. Array handling

Array handling is included in Fortran 90 for two main reasons:

- the notational convenience it provides, bringing the code closer to the underlying mathematical form;
- for the additional optimization opportunities it gives compilers (although there are plenty of opportunities for degrading optimization too!).

At the same time, major extensions of the functionality in this area have been added.

We have already met whole arrays in Parts 1 and 2—here we develop the theme.

Zero-sized arrays

A zero-sized array is handled by Fortran 90 as a legitimate object, without special coding by the programmer. Thus, in

```
DO i = 1,n
   x(i) = b(i) / a(i, i)
   b(i+1:n) = b(i+1:n) - a(i+1:n, i) * x(i)
END DO
```

no special code is required for the final iteration where i = n.

We note that a zero-sized array is regarded as being defined; however, an array of shape, say, (0,2) is not conformable with one of shape (0,3), whereas

```
x(1:0) = 3
```

is a valid “do nothing” statement.

Assumed-shape arrays

These are an extension and replacement for assumed-size arrays. Given an actual argument like:

```
REAL, DIMENSION(0:10, 0:20) :: a
   :
CALL sub(a)
```

the corresponding dummy argument specification defines only the type and rank of the array, not its size. This information has to be made available by an explicit interface, often using an interface block (see chapter 4). Thus we write just

```
SUBROUTINE sub(da)
   REAL, DIMENSION(:, :) :: da
```

and this is as if da were dimensioned (11,21). However, we can specify any lower bound and the array maps accordingly. The shape, not bounds, is passed, where the default lower bound is 1 and the defaultupper bound is the corresponding extent.

Automatic arrays

A partial replacement for the uses to which EQUIVALENCE is put is provided by this facility, useful for local, temporary arrays, as in

```
SUBROUTINE swap(a, b)
   REAL, DIMENSION(:)       :: a, b
   REAL, DIMENSION(SIZE(a)) :: work  ! array created on a stack
   work = a
   a = b
   b = work
END SUBROUTINE swap
```

ALLOCATABLE and ALLOCATE

Fortran 90 provides *dynamic* allocation of storage; it relies on a heap storage mechanism (and replaces another use of EQUIVALENCE). An example, for establishing a work array for a whole program, is

```
MODULE work_array
   INTEGER n
   REAL, DIMENSION(:,:,:), ALLOCATABLE :: work
END MODULE
PROGRAM main
   USE work_array
   READ (*, *) n
   ALLOCATE(work(n, 2*n, 3*n), STAT=status)
   :
   DEALLOCATE (work)
```

The work array can be propagated through the whole program via a USE statement in each program unit. We may specify an explicit lower bound and allocate several entities in one statement. To free dead storage we write, for instance,

```
DEALLOCATE(a, b)
```

We will meet this later, in the context of pointers.

Elemental operations and assignments

We have already met whole array assignments and operations:

```
REAL, DIMENSION(10) :: a, b
a = 0.          ! scalar broadcast; elemental assignment
b = sqrt(a)     ! intrinsic function result as array object
```

In the second assignment, an intrinsic function returns an array-valued result for an array-valued argument. We can write array-valued functions ourselves (they require an explicit interface):

```
PROGRAM test
   REAL, DIMENSION(3) :: a = (/ 1., 2., 3./),  b = (/ 2., 2., 2. /),  r
   r = f(a, b)
   PRINT *, r
CONTAINS
   FUNCTION f(c, d)
   REAL, DIMENSION(:) :: c, d
   REAL, DIMENSION(SIZE(c)) :: f
   f = c*d        ! (or some more useful function of c and d)
   END FUNCTION f
END PROGRAM test
```

WHERE

Often, we need to mask an assignment. This we can do using the WHERE, either as a statement:

```
WHERE (a /= 0.0) a = 1.0/a  ! avoid division by 0
```

(note: test is element-by-element, not on whole array), or as a construct (all arrays of same shape):

```
WHERE (a /= 0.0)
   a = 1.0/a
   b = a
END WHERE
```

```
WHERE (a /= 0.0)
   a = 1.0/a
ELSEWHERE
   a = HUGE(a)
END WHERE
```

Array elements

Simple case: given `REAL, DIMENSION(100, 100) :: a`
we can reference a single element of `a` as, for instance, `a(1, 1)`. For a derived data type like

```
TYPE triplet
   REAL                  u
   REAL, DIMENSION(3) :: du
END TYPE triplet
```

we can declare an array of that type: `TYPE(triplet), DIMENSION(10, 20) :: tar`
and a reference like

```
tar(n, 2)
```

is an element (a scalar!) of type triplet, but

```
tar(n, 2)%du
```

is an array of type real, and

```
tar(n, 2)%du(2)
```

is an element of it. The basic rule to remember is that an array element always has a subscript or subscripts qualifying at least the last name.

Array subobjects (sections)

The general form of subscript for an array section is

```
[lower] : [upper] [:stride]
```

as in

```
REAL a(10, 10)
a(i, 1:n)                 ! part of one row
a(1:m, j)                 ! part of one column
a(i, : )                  ! whole row
a(i, 1:n:3)               ! every third element of row
a(i, 10:1:-1)             ! row in reverse order
a( (/ 1, 7, 3, 2 /), 1)  ! vector subscript
a(1, 2:11:2)              ! 11 is legal as not referenced
a(:, 1:7)                 ! rank two section
```

Note that a vector subscript with duplicate values cannot appear on the left-hand side of an assignment as it would be ambiguous. Thus,

```
b( (/ 1, 7, 3, 7 /) ) = (/ 1, 2, 3, 4 /)
```

is illegal. Also, a section with a vector subscript must not be supplied as an actual argument to an `OUT` or `INOUT` dummy argument.

Arrays of arrays are not allowed:

```
tar%du             ! illegal
```

We note that a given value in an array can be referenced both as an element and as a section:

```
a(1, 1)           !  scalar (rank zero)
a(1:1, 1)         !  array section (rank one)
```

depending on the circumstances or requirements.
By qualifying objects of derived type, we obtain elements or sections depending on the rule stated earlier:

```
tar%u             !  array section (structure component)
tar(1, 1)%u       !  component of an array element
```

Arrays intrinsic functions

Vector and matrix multiply

```
DOT_PRODUCT       Dot product of 2 rank-one arrays
MATMUL            Matrix multiplication
```

Array reduction

```
ALL               True if all values are true
ANY               True if any value is true. Example: IF (ANY( a > b)) THEN
COUNT             Number of true elements in array
MAXVAL            Maximum value in an array
MINVAL            Minimum value in an array
PRODUCT           Product of array elements
SUM               Sum of array elements
```

Array inquiry

```
ALLOCATED         Array allocation status
LBOUND            Lower dimension bounds of an array
SHAPE             Shape of an array (or scalar)
SIZE              Total number of elements in an array
UBOUND            Upper dimension bounds of an array
```

Array construction

```
MERGE             Merge under mask
PACK              Pack an array into an array of rank
SPREAD            Replicate array by adding a dimension
UNPACK            Unpack an array of rank one into an array under mask
```

Array reshape

```
RESHAPE           Reshape an array
```

Array manipulation

```
CSHIFT            Circular shift
EOSHIFT           End-off shift
TRANSPOSE         Transpose of an array of rank two
```

Array location

```
MAXLOC            Location of first maximum value in an array
MINLOC            Location of first minimum value in an array
```

6. Pointers

Basics

Pointers are variables with the `POINTER` attribute; they are not a distinct data type (and so no "pointer arithmetic" is possible):

```
REAL, POINTER :: var
```

They are conceptually a descriptor listing the attributes of the objects (targets) that the pointer may point to, and the address, if any, of a target. They have no associated storage until it is allocated or otherwise associated (by pointer assignment, see below):

```
ALLOCATE (var)
```

and they are dereferenced automatically, so no special symbol is required. In

```
var = var + 2.3
```

the value of the target of `var` is used and modified. Pointers cannot be transferred via I/O—the statement

```
WRITE *, var
```

writes the value of the target of `var` and not the pointer descriptor itself.

A pointer can point to other pointers, and hence to their targets, or to a static object that has the TARGET attribute:

```
REAL, POINTER :: object
REAL, TARGET  :: target_obj
var => object                  ! pointer assignment
var => target_obj
```

but they are strongly typed:

```
INTEGER, POINTER :: int_var
var => int_var                 ! illegal - types must match
```

and, similarly, for arrays the ranks as well as the type must agree.

A pointer can be a component of a derived data type:

```
TYPE entry                       ! type for sparse matrix
   REAL value
   INTEGER index
   TYPE(entry), POINTER :: next  ! note recursion
END TYPE entry
```

and we can define the beginning of a linked chain of such entries:

```
TYPE(entry), POINTER :: chain
```

After suitable allocations and definitions, the first two entries could be addressed as

```
chain%value           chain%next%value
chain%index           chain%next%index
chain%next            chain%next%next
```

but we would normally define additional pointers to point at, for instance, the first and current entries in the list.

Association

A pointer's association status is one of

- undefined (initial state);
- associated (after allocation or a pointer assignment);
- disassociated:

```
DEALLOCATE (p, q)  ! for returning storage
NULLIFY (p, q)     ! for setting to 'null'
```

Some care has to be taken not to leave a pointer "dangling" by use of `DEALLOCATE` on its target without `NULLIFY`ing any other pointer referring to it.

The intrinsic function `ASSOCIATED` can test the association status of a defined pointer:

```
IF (ASSOCIATED(pointer)) THEN
```

or between a defined pointer and a defined target (which may, itself, be a pointer):

```
IF (ASSOCIATED(pointer, target)) THEN
```

Pointers in expressions and assignments

For intrinsic types we can "sweep" pointers over different sets of target data using the same code without any data movement. Given the matrix manipulation `y = B C z`, we can write the following code (although, in this case, the same result could be achieved more simply by other means):

```
REAL, TARGET  :: b(10,10), c(10,10), r(10), s(10, z(10)
REAL, POINTER :: a(:,:), x(:), y(:)
INTEGER mult
:
DO mult = 1, 2
   IF (mult == 1) THEN
      y => r              ! no data movement
      a => c
      x => z
   ELSE
      y => s              ! no data movement
      a => b
      x => r
   END IF
   y = MATMUL(a, x)       ! common calculation
END DO
```

For objects of derived data type we have to distinguish between pointer and normal assignment. In

```
TYPE(entry), POINTER :: first, current
:
first => current
```

the assignment causes first to point at current, whereas

```
first =  current
```

causes current to overwrite first and is equivalent to

```
first%value = current%value
first%index = current%index
first%next => current%next
```

Pointer arguments

If an actual argument is a pointer then, if the dummy argument is also a pointer,

- it must have same rank,
- it receives its association status from the actual argument,
- it returns its final association status to the actual argument (note: the target may be undefined!),
- it may not have the INTENT attribute (it would be ambiguous),
- it requires an interface block.

If the dummy argument is not a pointer, it becomes associated with the target of the actual argument:

```
   REAL, POINTER :: a(:,:)
      :
   ALLOCATE (a(80, 80))
      :
   CALL sub(a)
      :
SUBROUTINE sub(c)
   REAL c(:, :)
```

Pointer functions

Function results may also have the `POINTER` attribute; this is useful if the result size depends on calculations performed in the function, as in

```
   USE data_handler
   REAL x(100)
   REAL, POINTER :: y(:)
   :
   y => compact(x)
```

where the module `data_handler` contains

```
   FUNCTION compact(x)
      REAL, POINTER :: compact(:)
      REAL x(:)
! A procedure to remove duplicates from the array x
      INTEGER n
      :              ! Find the number of distinct values, n
      ALLOCATE(compact(n))
      :              ! Copy the distinct values into compact
   END FUNCTION compact
```

The result can be used in an expression (but must be associated with a defined target).

Arrays of pointers

These do not exist as such: given

```
   TYPE(entry) :: rows(n)
```

then

```
   rows%next              ! illegal
```

would be such an object, but with an irregular storage pattern. For this reason they are not allowed. However, we can achieve the same effect by defining a derived data type with a pointer as its sole component:

```
TYPE row
   REAL, POINTER :: r(:)
END TYPE
```

and then defining arrays of this data type:

```
TYPE(row) :: s(n), t(n)
```

where the storage for the rows can be allocated by, for instance,

```
DO i = 1, n
   ALLOCATE (t(i)%r(1:i)) ! Allocate row i of length i
END DO
```

The array assignment

```
s = t
```

is then equivalent to the pointer assignments

```
s(i)%r => t(i)%r
```

for all components.

Pointers as dynamic aliases

Given an array

```
REAL, TARGET :: table(100,100)
```

that is frequently referenced with the fixed subscripts

```
table(m:n, p:q)
```

these references may be replaced by

```
REAL, DIMENSION(:, :), POINTER :: window
   :
window => table(m:n, p:q)
```

The subscripts of window are `1:n-m+1`, `1:q-p+1`. Similarly, for

```
tar%u
```

(as defined in chapter 5, page 15), we can use, say,

```
taru => tar%u
```

to point at all the u components of tar, and subscript it as

```
taru(1, 2)
```

The subscripts are as those of tar itself. (This replaces yet more of `EQUIVALENCE`.)

The source code of an extended example of the use of pointers to support a data structure can be obtained by anonymous ftp to `jkr.cc.rl.ac.uk` (130.246.8.23). The directory is `/pub/MandR` and the file name is `appxg.f90`.

7. Specification Statements

This part completes what we have learned so far about specification statements.

Implicit typing

The implicit typing rules of Fortran 77 still hold. However, it is good practice to explicitly type all variables, and this can be forced by inserting the statement

```
IMPLICIT NONE
```

at the beginning of each prorgam unit.

PARAMETER attribute

A named constant can be specified directly by adding the **PARAMETER** attribute and the constant values to a type statement:

```
REAL, DIMENSION(3), PARAMETER :: field = (/ 0., 1., 2. /)
TYPE(triplet), PARAMETER      :: t = triplet( 0., (/ 0., 0., 0. /) )
```

DATA statement

The **DATA** statement can be used also for arrays and variables of derived type. It is also the only way to initialise just parts of such objects, as well as to initialise to binary, octal or hexadecimal values:

```
TYPE(triplet) :: t1, t2
DATA t1/triplet( 0., (/ 0., 1., 2. /) )/, t2%u/0./ ! only one component of t2 initialized
DATA array(1:64) / 64*0/                           ! only a section of array initialized
DATA i, j, k/ B'01010101', O'77', Z'ff'/
```

Characters

There are many variations on the way character arrays may be specified. Among the shortest and longest are

```
CHARACTER name(4, 5)*20
CHARACTER (KIND = kanji, LEN = 20), DIMENSION (4, 5) :: name
```

Initialization expressions

The values used in **DATA** and **PARAMETER** statements, or in specification statements with these attributes, are constant expressions that may include references to: array and structure constructors, elemental intrinsic functions with integer or character arguments and results, and the six transformational functions **REPEAT**, **SELECTED_INT_KIND**, **TRIM**, **SELECTED_REAL_KIND**, **RESHAPE** and **TRANSFER**:

```
INTEGER, PARAMETER :: long = SELECTED_REAL_KIND(12), array(3) = (/ 1, 2, 3 /)
```

Specification expressions

It is possible to specify details of variables using any non-constant, scalar, integer expression that may also include inquiry function references:

```
SUBROUTINE s(b, m, c)
  USE mod                                 ! contains a
  REAL, DIMENSION(:, :)              :: b   ! assumed-shape array
  REAL, DIMENSION(UBOUND(b, 1) + 5)  :: x   ! automatic array
  INTEGER                               m
  CHARACTER(LEN=*)                      c   ! assumed-length
  CHARACTER(LEN= m + LEN(c))            cc  ! automatic object
  REAL (SELECTED_REAL_KIND(2*PRECISION(a))) z   ! precision of z twice that of a
```

PUBLIC and PRIVATE

These attributes are used in specifications in modules to limit the scope of entities. The attribute form is

```
REAL, PUBLIC     :: x, y, z           ! default
INTEGER, PRIVATE :: u, v, w
```

and the statement form is

```
PUBLIC  :: x, y, z, OPERATOR(.add.)
PRIVATE :: u, v, w, ASSIGNMENT(=), OPERATOR(*)
```

The statement form has to be used to limit access to operators, and can also be used to change the overall default:

```
PRIVATE                               ! sets default for module
PUBLIC  :: only_this
```

For a derived data type there are three possibilities: the type and its components are all `PUBLIC`, the type is `PUBLIC` and its components `PRIVATE` (the type only is visible and one can change its details easily), or all of it is `PRIVATE` (for internal use in the module only):

```
MODULE mine
   PRIVATE
   TYPE, PUBLIC :: list
      REAL x, y
      TYPE(list), POINTER :: next
   END TYPE list
   TYPE(list) :: tree
   :
END MODULE mine
```

USE statement

To gain access to entities in a module, we use the `USE` statement. It has options to resolve name clashes if an imported name is the same as a local one:

```
USE mine, local_list => list
```

or to restrict the used entities to a specified set:

```
USE mine, ONLY : list
```

8. Intrinsic Procedures

We have already met most of the new intrinsic functions in previous parts of this series. Here, we deal only with their general classification and with those that have so far been omitted.

All intrinsic procedures can be referenced using keyword arguments:

```
CALL DATE_AND_TIME (TIME=t)
```

and many have optional arguments. They are grouped into four categories:

1. elemental – work on scalars or arrays, e.g. `ABS(a)`;
2. inquiry – independent of value of argument (which maybe undefined), e.g. `PRECISION(a)`;
3. transformational – array argument with array result of different shape, e.g. `RESHAPE(a, b)`;
4. subroutines, e.g. `SYSTEM_CLOCK`.

The procedures not already introduced are:

- Bit inquiry

```
BIT_SIZE           Number of bits in the model
```

- Bit manipulation

```
BTEST              Bit testing
IAND               Logical AND
IBCLR              Clear bit
IBITS              Bit extraction
IBSET              Set bit
IEOR               Exclusive OR
IOR                Inclusive OR
ISHFT              Logical shift
ISHFTC             Circular shift
NOT                Logical complement
```

- Transfer function, as in

```
INTEGER :: i = TRANSFER('abcd', 0) ! replaces part of EQUIVALENCE
```

- Subroutines

```
DATE_AND_TIME      Obtain date and/or time
MVBITS             Copies bits
RANDOM_NUMBER      Returns pseudorandom numbers
RANDOM_SEED        Access to seed
SYSTEM_CLOCK       Access to system clock
```

9. Input/Output

Non-advancing input/output

Normally, records of external, formatted files are positioned at their ends after a read or write operation. This can now be overridden with the additional specifiers:

```
  ADVANCE = 'NO'            (default is 'YES')
  EOR = eor_label           (optional, READ only)
  SIZE = size               (optional, READ only)
```

The next example shows how to read a record three characters at a time, and to take action if there are fewer than three left in the record:

```
     CHARACTER(3) key
     INTEGER unit, size
     READ (unit, '(A3)', ADVANCE='NO', SIZE=size, EOR=66) key
     :
! key is not in one record
  66 key(size+1:) = ''
     :
```

This shows how to keep the cursor positioned after a prompt:

```
     WRITE (*, '(A)', ADVANCE='NO') 'Enter next prime number:'
     READ  (*, '(I10)')          prime_number
```

New edit descriptors

The first three new edit descriptors are modelled on the `I` edit descriptor:

`B` binary,
`O` octal,
`Z` hexadecimal.

There are two new descriptors for real numbers:

`EN` engineering, multiple-of-three exponent: `0.0217  -->  21.70E-03  (EN9.2)`
`ES` scientific, leading nonzero digit: `0.0217  -->  2.17E-02  (ES9.2)`

and the `G` edit descriptor is generalized to all intrinsic types (`E/F`, `I`, `L`, `A`).

For entities of derived types, the programmer must elaborate a format for the ultimate components:

```
     TYPE string
        INTEGER length
        CHARACTER(LEN=20) word
     END TYPE string
     TYPE(string) :: text
     READ(*, '(I2, A)') text
```

New specifiers

On the `OPEN` and `INQUIRE` statements there are new specifiers:

```
POSITION =    'ASIS'       'REWIND'   'APPEND'
ACTION   =    'READ'       'WRITE'    'READWRITE'
DELIM    =    'APOSTROPHE' 'QUOTE'    'NONE'
PAD      =    'YES'        'NO'
```

and on the `INQUIRE` there are also

```
READ     = )
WRITE    = )  'YES'        'NO'       'UNKNOWN'
READWRITE= )
```

Finally, inquiry by I/O list (unformatted only) is possible:

```
INQUIRE (IOLENGTH = length) item1, item2,...
```

and this is useful to set `RECL`, or to check that a list is not too long. It is in the same processor-dependent units as `RECL` and thus is a portability aid.

Index